

EULERIAN AND HAMILTONIAN PROPERTIES OF GALLAI AND ANTI-GALLAI TOTAL GRAPHS

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Abstract. Let $G = (V, E)$ be a graph. The *Gallai total graph* $\Gamma_T(G)$ of G is the graph, where $V(\Gamma_T(G)) = V \cup E$ and $uv \in E(\Gamma_T(G))$ if and only if

- (i) u and v are adjacent vertices in G , or
- (ii) u is incident to v or v is incident to u in G , or
- (iii) u and v are adjacent edges in G which do not span a triangle in G .

The *anti-Gallai total graph* $\Delta_T(G)$ of G is the graph, where $V(\Delta_T(G)) = V \cup E$ and $uv \in E(\Delta_T(G))$ if and only if

- (i) u and v are adjacent vertices in G , or
- (ii) u is incident to v or v is incident to u in G , or
- (iii) u and v are adjacent edges in G and lie on a same triangle in G .

In this paper, we discuss Eulerian and Hamiltonian properties of Gallai and anti-Gallai total graphs.

Key words: Euler graph, Hamiltonian graph, Gallai total graph, anti-Gallai total graph.

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Abstrak. Misalkan $G = (V, E)$ adalah sebuah graf. *Graf total Gallai* $\Gamma_T(G)$ dari G adalah sebuah graf dimana $V(\Gamma_T(G)) = V \cup E$ dan $uv \in E(\Gamma_T(G))$ jika dan hanya jika

- (i) u dan v adalah titik-titik bertetangga di G , atau
- (ii) u berkaitan dengan v atau v berkaitan dengan u di G , atau
- (iii) u dan v adalah sisi-sisi bertetangga di G yang tidak membangun suatu segitiga di G .

Graf total anti-Gallai $\Delta_T(G)$ dari G adalah sebuah graf dimana $V(\Delta_T(G)) = V \cup E$ dan $uv \in E(\Delta_T(G))$ jika dan hanya jika

- (i) u dan v adalah titik-titik bertetangga di G , atau
- (ii) u berkaitan dengan v atau v berkaitan dengan u di G , atau
- (iii) u dan v adalah sisi-sisi bertetangga di G dan terletak pada suatu segitiga di G .

Pada paper ini, kami meneliti sifat-sifat Euler and Hamilton dari graf total Gallai and anti-Gallai.

Kata kunci: Graf Euler, graf Hamilton, graf total Gallai, graf total anti-Gallai.

1. INTRODUCTION

A *graph* $G = (V, E)$ is an ordered pair of set of vertices and edges, where edges are unordered pair of vertices. Also G is said to be a (p, q) graph if $|V| = p$ and $|E| = q$. A graph is *simple* if it has neither self loop nor multiple edges. A *finite graph* is a graph with finite number of vertices and edges. Two vertices (edges) are said to be *adjacent* if they have a common edge (vertex). If a vertex v lies on an edge e , then they are said to be *incident* to each other. The *degree* $d(v)$ of a vertex $v \in V$ is the number of edges incident at v . A *regular graph* is the graph in which every vertex of the graph has same degree. A *path* in a graph G is a walk with no repeated vertex. A graph G is said to be *connected* if there exists a path between every pair of vertices of G . Let $G = (V, E)$ be a graph with $|V| = p$, then the *adjacency matrix* $A(G)$ of G is defined as $A(G) = [a_{ij}]_{p \times p}$, where

$$a_{ij} = \begin{cases} 1 & \text{if } v_i \text{ is adjacent to } v_j, i \neq j, \\ 0 & \text{otherwise.} \end{cases}$$

A graph G is called *Euler graph* if there exists a closed walk in G with no repeated edge and all the edges are traversed exactly once. A closed path is called a *cycle*. A cycle is said to be *spanning cycle* if it contains all the vertices of the graph. A graph G is said to be *Hamiltonian* if it contains a spanning cycle. Vertices and edges of G are called *elements* of G .

The *line graph* $L(G)$ of a graph G is defined as the graph whose vertices are the edges of G , with two vertices adjacent in $L(G)$ if and only if the corresponding edges are adjacent in G . The line graphs were first studied by Whitney [21]. The line graph is the well studied graph operator in the literature (see [7], [8], [19]).

The *Gallai graph* $\Gamma(G)$ of a graph G is the graph in which $V(\Gamma(G)) = E(G)$ and two distinct edges of G are adjacent in $\Gamma(G)$ if they are adjacent in G , but do not span a triangle in G . The *anti-Gallai graph* $\Delta(G)$ of a graph G is the graph in which $V(\Delta(G)) = E(G)$ and two distinct edges of G are adjacent in $\Delta(G)$ if they are adjacent in G and lie on a same triangle in G . These constructions were used by Gallai [16] in his investigation of comparability graphs; the notion was suggested by Sun [20]. Sun used the Gallai graphs to describe a nice class of perfect graphs. Gallai graphs are also used in polynomial time algorithm to recognize $k_{1,3}$ -free perfect graphs by Chvatal and Sbihi [15]. Several properties of Gallai and anti-Gallai graphs are discussed in [1], [2], [3], [4], [5], [6] and [18].

Remark: A graph is *triangle free* if it does not contains a triangle. If a graph G is triangle free, then $\Gamma(G) \cong L(G)$.

The *total graph* $T(G)$ of G is the graph whose vertex set is $V(G) \cup E(G)$, and two vertices are adjacent if and only if they are adjacent or incident in G . The notion of total graph was introduced by Behzad and Chartrand [12]. Several properties of total graphs are investigated in the literature (see [9], [10], [13], [14]). Behzad obtained a characterization of total graphs in [11]. Gavril established a linear time algorithm for the recognition of the total graphs in [17].

Motivated from the operators Gallai graphs, anti-Gallai graphs and total graphs, we introduce two operators Gallai and anti-Gallai total graphs as follows: Let $G = (V, E)$ be a graph. The *Gallai total graph* $\Gamma_T(G)$ of G is the graph, where $V(\Gamma_T(G)) = V \cup E$ and $uv \in E(\Gamma_T(G))$ if and only if

- (i) u and v are adjacent vertices in G , or
- (ii) u is incident to v or v is incident to u in G , or
- (iii) u and v are adjacent edges in G which do not span a triangle in G .

The *anti-Gallai total graph* $\Delta_T(G)$ of G is the graph, where $V(\Delta_T(G)) = V \cup E$ and $uv \in E(\Delta_T(G))$ if and only if

- (i) u and v are adjacent vertices in G , or
- (ii) u is incident to v or v is incident to u in G , or
- (iii) u and v are adjacent edges in G and lie on a same triangle in G .

The Gallai total graph $H = \Gamma_T(G)$ and anti-Gallai total graph $H' = \Delta_T(G)$ of G are shown in Figure 1. Degree of $v \in V(\Gamma_T(G))$ and $v \in V(\Delta_T(G))$ are denoted by $d_\Gamma(v)$ and $d_\Delta(v)$, respectively. Throughout the paper, we consider simple and finite graphs.

Remark: If the graph G is triangle free, then $\Gamma_T(G) \cong T(G)$.

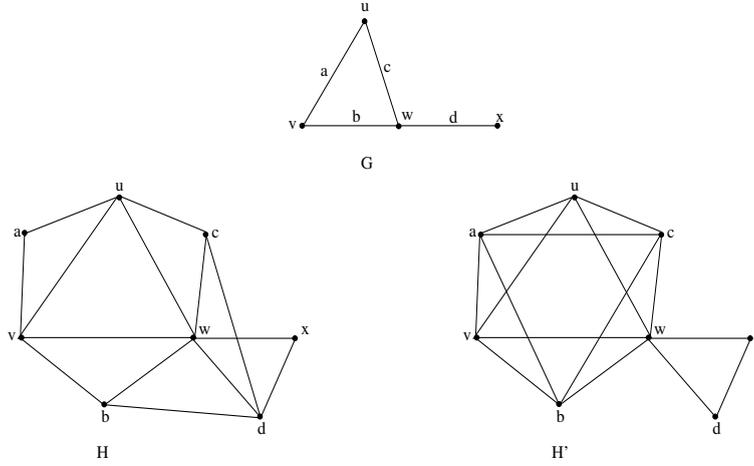


FIGURE 1. A graph G , its Gallai total graph $H = \Gamma_T(G)$ and anti-Gallai total graph $H' = \Delta_T(G)$

2. EULERIAN GALLAI TOTAL GRAPHS

Proposition 2.1. *Let $\Gamma_T(G)$ be the Gallai total graph of a graph $G = (V, E)$, then*

$$d_{\Gamma}(v) = \begin{cases} 2d(v); & \text{if } v \in V(G), \\ d(v_1) + d(v_2) - 2t; & \text{if } v = v_1v_2 \in E(G), \end{cases}$$

where t denotes the number of triangles containing v in G .

Proof. Let $\Gamma_T(G)$ be the corresponding Gallai total graph of G . Since the graph G is a subgraph of $\Gamma_T(G)$ and also each edge incident to v in G is adjacent to corresponding vertex v in $\Gamma_T(G)$, therefore, $d_{\Gamma}(v) = 2d(v)$, for $v \in V(G)$. If $v = v_1v_2 \in E(G)$, then the corresponding vertex v of $\Gamma_T(G)$ is adjacent to all the edges which are adjacent to v , but do not lie on a same triangle with v in G . It implies that they contribute the degree $(d(v_1) - 1) + (d(v_2) - 1) - 2t$ (if v is the edge of a triangle, then it is not adjacent to those two edges of G which span a triangle with v in G) and v is also adjacent in $\Gamma_T(G)$ to the vertices by which it is incident in G , so 2 more degrees are also contributed, therefore, $d_{\Gamma}(v) = d(v_1) + d(v_2) - 2t$. \square

Lemma 2.2. *The number of triangles in a graph G is equal to $\frac{tr(A^3)}{6}$, where A is the adjacency matrix of G .*

Proposition 2.3. *Let $G = (V, E)$ be a (p, q) graph, then $|E(\Gamma_T(G))| = 2q + \frac{1}{2} \sum_{i=1}^p (d(v_i))^2 - 3 \left(\frac{tr(A^3)}{6} \right)$, where A is the adjacency matrix of G .*

Proof. Let $G = (V, E)$ be a (p, q) graph and $v_1, v_2, \dots, v_i, \dots, v_p$ be vertices of G . Then total degree of vertices of $\Gamma_T(G)$ is equal to 2 (sum of degree of the vertices of

G) + sum of the degree of the vertices corresponding to the edges of G . Let $E'(G)$ be the set of edges which do not lie on a triangle in G and $|E'(G)| = q_1$. Also let $E''(G)$ be the set of edges which lie on a triangle in G and $|E''(G)| = q_2$. Now if $e = v_i v_j \in E'(G)$, then degree of the corresponding vertex e' in $\Gamma_T(G)$ is equal to $d(v_i) + d(v_j)$, so the total degree of the vertices in $\Gamma_T(G)$ corresponding to such edges of G is $\sum_{v_i v_j \in E'(G)}^{q_1} (d(v_i) + d(v_j))$. If $e = v_i v_j \in E''(G)$, then degree of the corresponding vertex e' in $\Gamma_T(G)$ is equal to $d(v_i) + d(v_j) - 2t_{ij}$, where t_{ij} is the number of triangles on which the edge $v_i v_j$ lies, so the total degree of the vertices in $\Gamma_T(G)$ corresponding to such edges of G is $\sum_{v_i v_j \in E''(G)}^{q_2} (d(v_i) + d(v_j) - 2t_{ij})$. Then by handshake lemma on G and $\Gamma_T(G)$ we have, total degree of $\Gamma_T(G)$

$$\begin{aligned}
&= 2(2q) + \sum_{v_i v_j \in E'(G)}^{q_1} (d(v_i) + d(v_j)) + \sum_{v_i v_j \in E''(G)}^{q_2} (d(v_i) + d(v_j) - 2t_{ij}) \\
&= 4q + \sum_{v_i v_j \in E(G)}^q (d(v_i) + d(v_j)) - 2 \sum_{v_i v_j \in E''(G)}^{q_2} (t_{ij}) \\
&= 4q + \sum_{i=1}^p ((d(v_i))^2) - 2(3 \times \text{total no. of triangles in } G) \\
&= 4q + \sum_{i=1}^p (d(v_i))^2 - 6(\text{number of triangles in } G) \\
&= 4q + \sum_{i=1}^p (d(v_i))^2 - 6 \left(\frac{\text{tr}(A^3)}{6} \right), \text{ using Lemma 2.2.}
\end{aligned}$$

Then by handshake lemma, the total number of edges in $\Gamma_T(G)$,

$$|E(\Gamma_T(G))| = 2q + \frac{1}{2} \sum_{i=1}^p (d(v_i))^2 - 3 \left(\frac{\text{tr}(A^3)}{6} \right).$$

□

Proposition 2.4. *The Gallai total graph $\Gamma_T(G)$ of a graph G is regular if and only if G is regular and triangle free.*

Proof. Let G be a regular and triangle free graph. Now we have to show that $\Gamma_T(G)$ is regular. Since G is regular, degree of each vertex is same. It implies that corresponding vertices of $\Gamma_T(G)$ are of same degree (by the Proposition 2.1). Also it is given that G is triangle free. It follows that every vertex corresponding to the edges of G has degree $d(u) + d(v)$ (by the Proposition 2.1), where u and v are the end vertices of the edge and this is equal to $2d(u)$, being G is regular. Therefore, degree of each vertex of $\Gamma_T(G)$ is same. Hence, $\Gamma_T(G)$ is regular. Conversely, suppose that $\Gamma_T(G)$ is regular. Now we have to show that G is regular and triangle

free. Let on contrary G is neither regular nor triangle free. If G is not regular, then degree of every vertex of $\Gamma_T(G)$ corresponding to the vertices of G is not same, which is a contradiction to our fact that $\Gamma_T(G)$ is regular. Hence G is regular. Now if G is not triangle free, then degree of every vertex of $\Gamma_T(G)$ corresponding to the edges of G is $d(u) + d(v) - 2t$ (by the Proposition 2.1), where t is the number of triangles containing the edge, which is again a contradiction to our fact that $\Gamma_T(G)$ is regular. Thus, G is triangle free. Hence, G is regular and triangle free. \square

Proposition 2.5. *For a connected graph G , the Gallai total graph $\Gamma_T(G)$ has a spanning Eulerian subgraph.*

Proof. The result is obvious if G consists of a single vertex. Otherwise, the connected spanning subgraph H of $\Gamma_T(G)$ obtained by deleting all the edges of the Gallai graph $\Gamma(G)$, is Eulerian, since all its vertices have even degree. \square

Corollary 2.6. *For a connected graph G , $\Gamma_T^n(G)$ has a spanning Eulerian subgraph for all $n \geq 1$.*

Proposition 2.7. *The Gallai total graph $\Gamma_T(G)$ of G is connected if and only if G is connected.*

Proof. Necessity: Let G be a connected graph. By the Proposition 2.5, $\Gamma_T(G)$ has a spanning Eulerian subgraph. That means, there exists a path between every pair of vertices of $\Gamma_T(G)$. Hence, $\Gamma_T(G)$ is connected.

Sufficiency: Suppose $\Gamma_T(G)$ is connected. Now, we have to show that G is connected. Let on contrary G be disconnected. Then, there exists at least a pair of vertices which has no path between them. Let u, v be two such vertices, then there is no path between u and v in $\Gamma_T(G)$. It follows that $\Gamma_T(G)$ is disconnected, a contradiction to the hypothesis. Hence the theorem. \square

Corollary 2.8. *The $\Gamma_T^n(G)$ of G is connected if and only if G is connected for all $n \geq 1$.*

Parity of a vertex means parity of its degree, i.e. degree of the vertex is either even or odd.

Theorem 2.9. *Let G be a connected graph, then the Gallai total graph $\Gamma_T(G)$ is Eulerian if and only if all the vertices of G are of the same parity.*

Proof. Let $\Gamma_T(G)$ be an Eulerian graph. Assume to the contrary, G has a vertex v_1 with odd degree and a vertex v_2 with even degree. Since G is connected, v_1 and v_2 are joined by a path which contains two adjacent vertices v_i and v_j (not necessary different from v_1 or v_2) of opposite parity. Thus, by the Proposition 2.1 (part (ii)), the vertex in $\Gamma_T(G)$ which corresponds to the edge $v_i v_j$ has odd degree, a contradiction to the fact that $\Gamma_T(G)$ is Eulerian. Hence, all the vertices of G are

of the same parity.

Conversely, suppose all the vertices of G are of the same parity, then by the Proposition 2.1, the vertices of $\Gamma_T(G)$ are of even degree. Therefore, $\Gamma_T(G)$ is Eulerian. Hence the theorem. \square

Corollary 2.10. *If G is Eulerian, then $\Gamma_T^n(G)$ is Eulerian for all $n \geq 1$, but converse is not true.*

Counter example of converse: The Gallai total graph $H = \Gamma_T(G)$ of G is Eulerian, but G is not an Eulerian graph as shown in Figure 2.

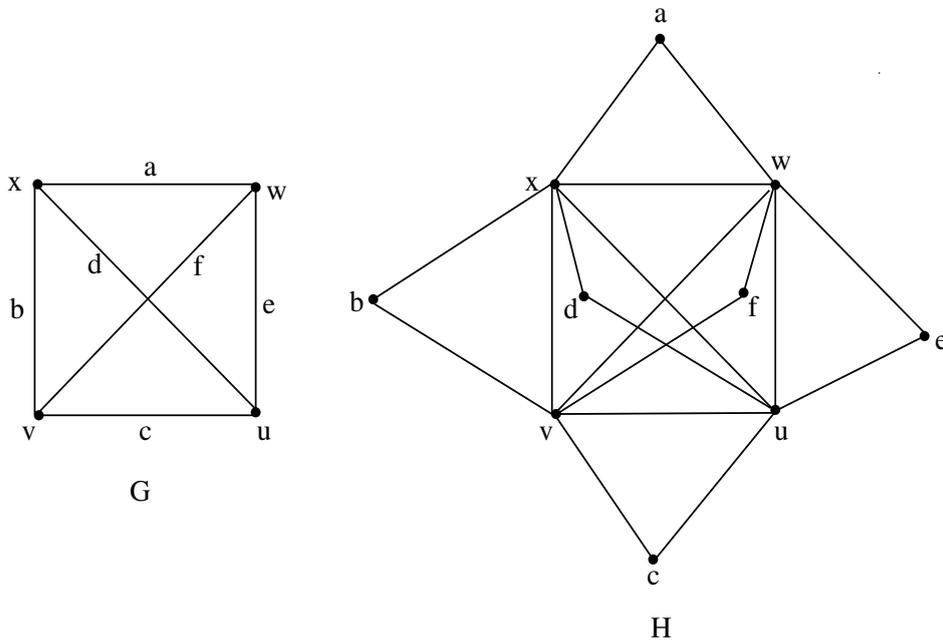


FIGURE 2. Eulerian Gallai total graph $H = \Gamma_T(G)$ of a non Eulerian graph G

3. EULERIAN ANTI-GALLAI TOTAL GRAPHS

Proposition 3.1. *Let $\Delta_T(G)$ be the anti-Gallai total graph of a graph $G = (V, E)$, then*

$$d_{\Delta}(v) = \begin{cases} 2d(v); & \text{if } v \in V(G), \\ 2t + 2; & \text{if } v = v_1v_2 \in E(G), \end{cases}$$

where t denotes the number of triangles containing v in G .

Proof. Let $\Delta_T(G)$ be the corresponding anti-Gallai total graph of G . Since the graph G is a subgraph of $\Delta_T(G)$ and also each edge incident to v in G is adjacent to corresponding vertex v in $\Delta_T(G)$, therefore, $d_{\Delta}(v) = 2d(v)$, for $v \in V(G)$. If $v = v_1v_2 \in E(G)$, then the corresponding vertex v of $\Delta_T(G)$ is adjacent to all the edges which are adjacent to v , and lie on a same triangle with v in G . It implies they contribute the degree $2t$ (if v is the edge of a triangle, then it is adjacent to those two edges of G which span triangle with v in G) and v is also adjacent in $\Delta_T(G)$ to the vertices by which it is incident in G , so 2 more degrees are also contributed, therefore, $d_{\Delta}(v) = 2t + 2$. □

Proposition 3.2. *Let $G = (V, E)$ be a (p, q) graph, then $|E(\Delta_T(G))| = 3q + 3\left(\frac{\text{tr}(A^3)}{6}\right)$, where A is the adjacency matrix of G .*

Proof. Let $G = (V, E)$ be a (p, q) graph and $v_1, v_2, \dots, v_i, \dots, v_p$ be vertices of G . Then by the definition of $\Delta_T(G)$, sum of the degree of its vertices is equal to 2(sum of degree of the vertices of G)+sum of the degree of the vertices corresponding to the edges of G . If the edge $e = v_iv_j$ does not lie on a triangle in G , then the degree of the corresponding vertex in $\Delta_T(G)$ is equal to 2, so the sum of the degree of the vertices corresponding to the edges of G is $2q$ (e is adjacent to vertices in $\Delta_T(G)$ corresponding to the vertices of G , by which it is incident). But if an edge lie on a triangle in G , then 2 degrees are increased, so if there is a triangle in G , then 6 degrees are increased. Therefore, sum of the degree of the vertices of $\Delta_T(G)$ is

$$\begin{aligned} &= 2(2q) + 2q + 6(\text{number of triangles in } G) \\ &= 4q + 2q + 6\left(\frac{\text{tr}(A^3)}{6}\right) \text{ (using Lemma 2.2)} \\ &= 6q + 6\left(\frac{\text{tr}(A^3)}{6}\right). \end{aligned}$$

By handshake lemma, the total number of edges in $\Delta_T(G)$,

$$|E(\Delta_T(G))| = 3q + 3\left(\frac{\text{tr}(A^3)}{6}\right).$$

□

A graph is called l -triangular if each edge of G lies on l number of triangles in G .

Proposition 3.3. *The anti-Gallai total graph $\Delta_T(G)$ of G is regular if and only if G is l -triangular and $(l + 1)$ -regular.*

Proof. Suppose G is l -triangular and $(l + 1)$ -regular. Now we have to show that $\Delta_T(G)$ is regular. Given that G is $(l + 1)$ -regular, therefore the degree of every vertex in $\Delta_T(G)$ corresponding to the vertices of G are $2(l + 1)$ (by the Proposition 3.1). Also G is l -triangular, It implies that the degree of every vertex of $\Delta_T(G)$ corresponding to the edges of G are $2l + 2$ (by the Proposition 3.1). Hence degree

of each vertex of $\Delta_T(G)$ is same. Therefore, $\Delta_T(G)$ is regular.

Conversely, suppose $\Delta_T(G)$ is regular. Now we have to show that the given conditions are satisfied. Let on contrary, G be neither l -triangular nor $(l + 1)$ -regular. If G is not $(l + 1)$ -regular, then the degree of each vertex in $\Delta_T(G)$ corresponding to the vertices of G is not same, which is a contradiction to our fact that $\Delta_T(G)$ is regular. Hence G is $(l + 1)$ -regular. Now, if G is not l -triangular, then the degree of the vertices corresponding to the edges of G are not same (by the Proposition 3.1), which is a contradiction to our fact that $\Delta_T(G)$ is regular. Thus, G is l -triangular. Hence G is l -triangular and $(l + 1)$ -regular. \square

Proposition 3.4. *For a connected graph G , the anti-Gallai total graph $\Delta_T(G)$ has a spanning Eulerian subgraph.*

Proof. The result is obvious if G consists of a single vertex. Otherwise, the connected spanning subgraph H of $\Delta_T(G)$ obtained by deleting all the edges of $\Delta(G)$, is Eulerian, since all its vertices have even degree. \square

Corollary 3.5. *For a connected graph G , $\Delta_T^n(G)$ has a spanning Eulerian subgraph for all $n \geq 1$.*

Proposition 3.6. *The anti-Gallai total graph $\Delta_T(G)$ of G is connected if and only if G is connected.*

Proof. Suppose G is a connected graph, then by the Proposition 3.4, $\Delta_T(G)$ has a spanning Eulerian subgraph. That means, there exists a path between every pair of vertices of $\Delta_T(G)$. Hence, $\Delta_T(G)$ is connected.

Conversely, suppose $\Delta_T(G)$ is connected. Now, we have to show that G is connected. Let on contrary, G be disconnected. Then, there exists at least a pair of vertices which has no path between them. Let u, v be two such vertices, then by the definition of $\Delta_T(G)$, u and v has no path in $\Delta_T(G)$. It follows that $\Delta_T(G)$ is disconnected graph, a contradiction to the hypothesis. Hence the theorem. \square

Corollary 3.7. *The anti-Gallai total graph $\Delta_T^n(G)$ of G is connected if and only if G is connected for all $n \geq 1$.*

Theorem 3.8. *Let G be a connected graph, then $\Delta_T(G)$ of G is Eulerian.*

Proof. By the Proposition 3.1, all the vertices of $\Delta_T(G)$ are of even degree and given that G is connected. That means, $\Delta_T(G)$ is Eulerian. \square

Corollary 3.9. *Let G be a connected graph, then $\Delta_T^n(G)$ of G is Eulerian.*

Corollary 3.10. *If G is Eulerian, then $\Delta_T^n(G)$ of G is Eulerian, but converse is not true.*

Counter example of converse: The anti-Gallai total graph $\Delta_T(G)$ of G is Eulerian, but G is not an Eulerian graph as shown in Figure 3.

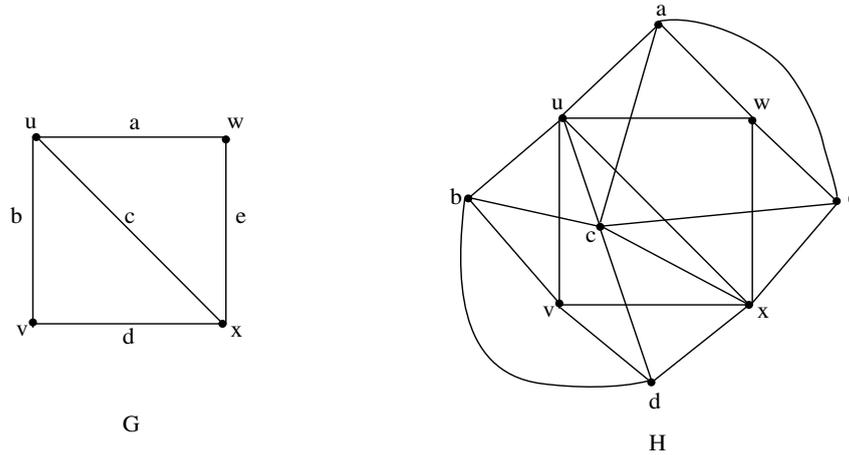


FIGURE 3. Eulerian anti-Gallai total graph $H = \Delta_T(G)$ of a non Eulerian graph G

4. HAMILTONIAN GALLAI AND ANTI-GALLAI TOTAL GRAPHS

Theorem 4.1. *The Gallai total graph $\Gamma_T(G)$ of a non-trivial graph G is Hamiltonian if and only if the set of all elements of G can be ordered in such a way that consecutive elements are neighbour as are the first and last elements, but the two edges are not consecutive elements, if both the edges are of same triangle.*

Proof. Let G be a graph. The elements of G can be ordered in such a way that consecutive elements are neighbour as are the first and last elements, but the two edges are not consecutive elements, if both the edges are of same triangle. Then by taking the corresponding vertices of $\Gamma_T(G)$ in the same order, a Hamiltonian cycle is produced. Hence, $\Gamma_T(G)$ is Hamiltonian.

Conversely, let $\Gamma_T(G)$ be a Hamiltonian graph. It follows that it contains a Hamiltonian cycle,

$$C = (v_0, v_1, \dots, v_m = v_0).$$

Let a_i be that element of G associated with the vertex v_i . Thus, we get an ordering $a_0, a_1, \dots, a_{m-1}, a_m = a_0$ of elements of G . By the definition of $\Gamma_T(G)$, this ordering is in such a way that consecutive elements are neighbour as are the first and last elements, but the two edges are not consecutive elements, if both the edges are of same triangle. Hence the theorem. \square

Lemma 4.2. *If G is Hamiltonian, then its total graph $T(G)$ is also Hamiltonian.*

Corollary 4.3. *If G is Hamiltonian and triangle-free, then $\Gamma_T(G)$ is also Hamiltonian.*

Theorem 4.4. *The anti-Gallai total graph $\Delta_T(G)$ of a non-trivial graph G is Hamiltonian if and only if the set of all elements of G can be ordered in such a way that consecutive elements are neighbour as are the first and last elements, but the two edges are not consecutive elements, if both the edges are not of same triangle.*

Proof. Let G be a graph. The elements of G can be ordered in such a way that consecutive elements are neighbour as are the first and last elements, but the two edges are not consecutive elements, if both the edges are not of same triangle. Then by taking the corresponding vertices of $\Delta_T(G)$ of G in the same order, a Hamiltonian cycle is produced. Hence $\Delta_T(G)$ is Hamiltonian.

Conversely, let $\Delta_T(G)$ be a Hamiltonian graph. It follows that it contains a Hamiltonian cycle,

$$C = (v_0, v_1, \dots, v_m = v_0).$$

Let a_i be that element of G associated with the vertex v_i . Thus, we get an ordering $a_0, a_1, \dots, a_{m-1}, a_m = a_0$ of elements of G . By the definition of $\Delta_T(G)$, this ordering is in such a way that consecutive elements are neighbour as are the first and last elements, but the two edges are not consecutive elements, if both the edges are not of same triangle. Hence the theorem. \square

5. CONCLUDING REMARKS

In this note, we have introduced the notion of Gallai and anti-Gallai total graphs, and have investigated the Eulerian and Hamiltonian properties of these graph operators. Further, one can find other graph theoretic properties of these operators.

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