

Implementation of Pattern Matching Algorithm for Multimedia Files in Mail Function Detection

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Abstract— Now a days internet and mail based file transfer has increased enormously due to this server space required will be highly and also occurs largely. In existing system if we upload the same file which is present in the server also get uploaded and duplication occurs. We used a pattern matching algorithm it eliminate duplication and also to avoid time wastage in uploading the same file present in the server. During file upload pattern will be matched. If

pattern matched file won't be uploaded again it will simply matched the existing file it avoids uploading the file again. If pattern doesn't match it allow uploading the file. From this we save the memory space in the server and duplication doesn't occur.

Keywords—Feature Selection, Boyer-Moore Algorithm, Pattern Matching, Data mining.

I. INTRODUCTION

In the earlier years, design coordinating has been routinely utilized as a part of different PC applications, for instance, in editors, recovery of data (from content, picture, or sound), and seeking nucleotide or amino corrosive arrangement designs in genome and protein grouping databases. The present day design coordinating calculations coordinate the example precisely or around inside the content. A careful example coordinating is to discover every one of the events of a specific example ($x_1 x_2 \dots x_m$) of m -characters in a content ($y_1 y_2 \dots y_n$) of n -characters which are worked over a limited arrangement of characters of a letters in order set indicated by Σ and the extent of this set is equivalent to σ . The immediate route to this issue is to analyze the principal m -characters of the content and the example in some predefined request and, after a match or a crisscross, slide the whole example by one character in the forward bearing of the content. This procedure is rehased until the example is situated at the $(nm+1)$ position of the content.

This methodology is generally known as a beast power technique. To encourage this errand, a few calculations have been proposed, and these have their own particular focal points and restrictions in view of the example length, periodicity, and the sort of the content (for e.g., nucleotide or amino corrosive

successions or dialect characters, and so forth.). The vast majority of the surely understood calculations (see beneath for points of interest) work in two stages: i.e., the preprocessing stage and the pursuit stage. In the preprocessing stage, these calculations handle the example and utilize this data in the hunt stage to lessen the aggregate number of character examinations and subsequently diminish the general execution time. The effectiveness of a calculation predominantly relies on upon the inquiry stage. The primary goal behind the example coordinating calculations is to decrease the aggregate number of character correlations between the example and the content to build the general productivity. The change in the proficiency of an inquiry can be accomplished by adjusting the request in which the characters are looked at every endeavor and by picking a movement component that allows the skipping of a predefined number of characters in the content after every endeavor. Design coordinating calculations filter the content with the assistance of a window, whose size is equivalent to the length of the example. The initial step is to adjust the left closures of the window and the content and after that look at the relating characters of the window and the example; this method is known as endeavor. After a match or a bungle of the example, the content window is moved to one side. The inquiry is what number of characters are required to move the window on the content. This movement esteem fluctuates taking into

account the system utilized by different calculations. This technique is rehashed until the right end of the window.

II. EXISTING SYSTEM

In the existing system, if we upload the same file which is present in the server also get uploaded and duplication occurs.

III. PROPOSED SYSTEM

We used an algorithm, Boyer Moore which eliminates the duplication in uploading the same files present in server. During file upload, pattern matching is generated by the server. If the pattern of the file is matched with the server files pattern, then the file won't be uploaded again. If pattern of the file doesn't match, then it allows to upload the file.

MODULES: In our project we have two modules named User and Server.

USER

In the USER MODULE, client can transfer the records, get the documents, view and upgrade the profile, download their documents, forward their records and amid transfer she/he ask for the server for example coordinating.

SERVER

In the server module, during the file upload by the user, the server generate the pattern to the that file. And compare the pattern with the already existing files in the server. For this pattern matching, we used the boyer-moore algorithm in the server.

BOYER-MOORE ALGORITHM: The calculation contrasts the example P and the substring of arrangement T inside a sliding window in the privilege to-left request. The awful character standard and great addition guideline are utilized to decide the development of sliding window.

EXAMPLE

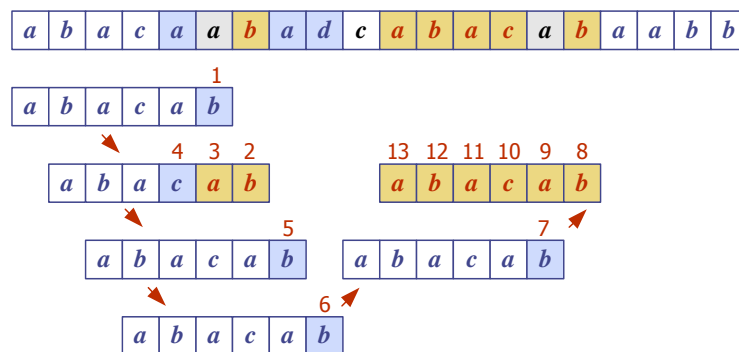
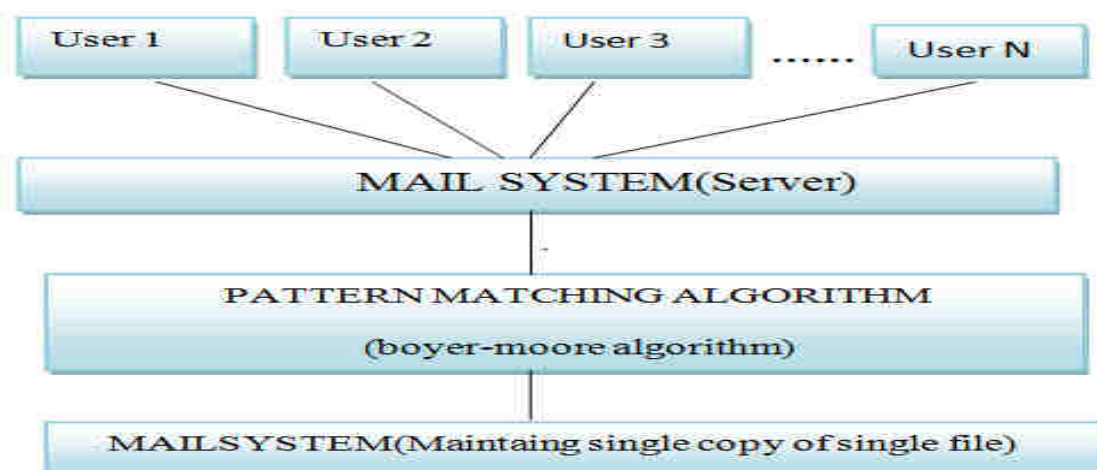


Fig.1: System Architecture

IV. SYSTEM ARCHITECTURE



V. RESULTS

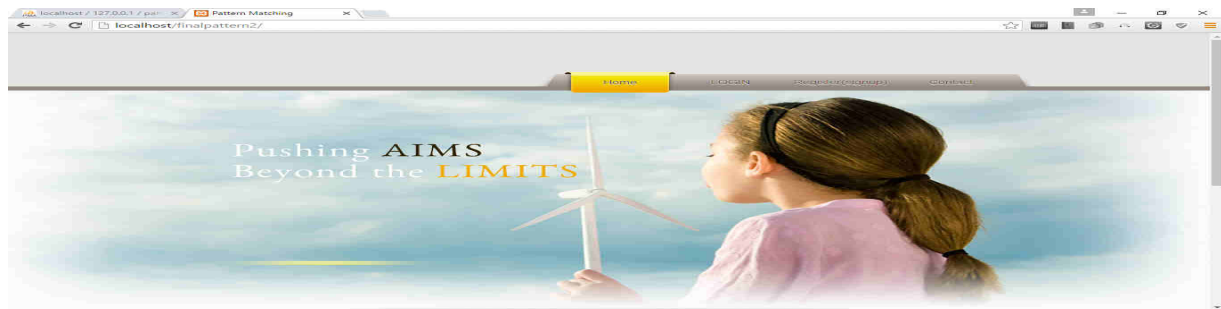


Fig. 2: HOME PAGE: This is our homepage for mail system



Fig. 3: HOME PAGE: This is our homepage for mail system

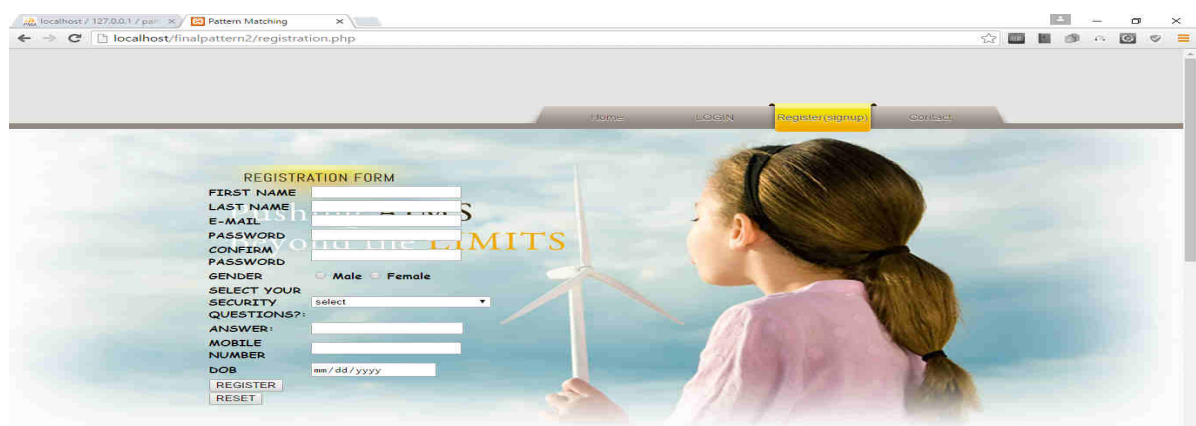


Fig.4: REGISTRATION PAGE: NEW USER can register into the mail system and create the new account in this mail system.



Fig.5: INBOX: In this user can view their mails.

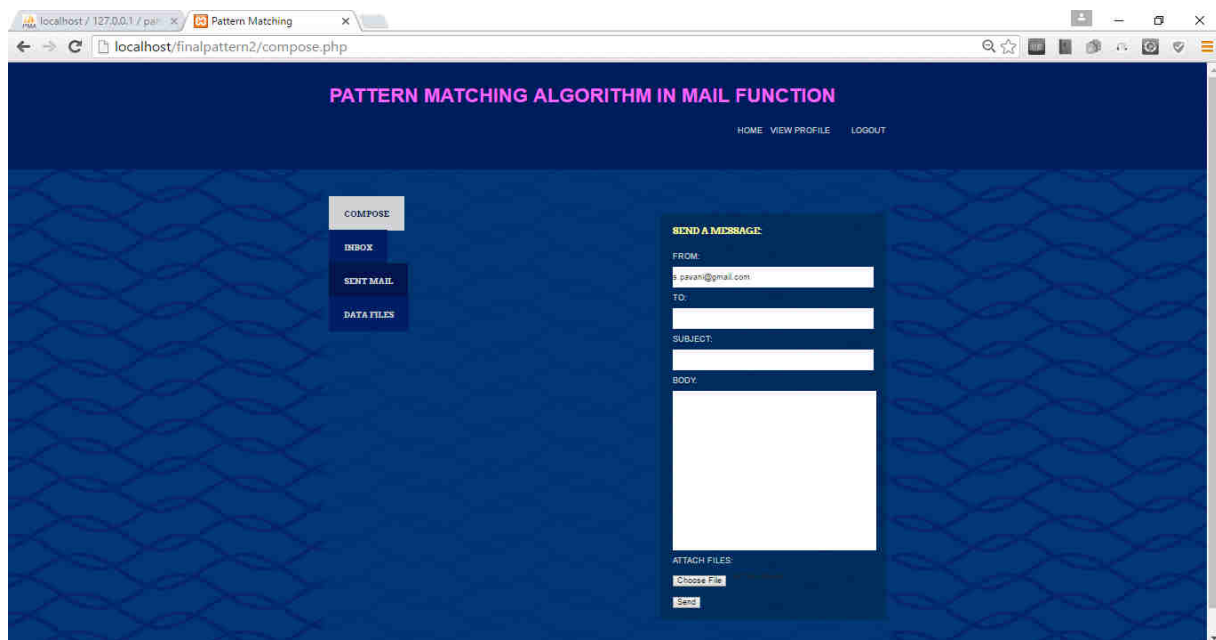
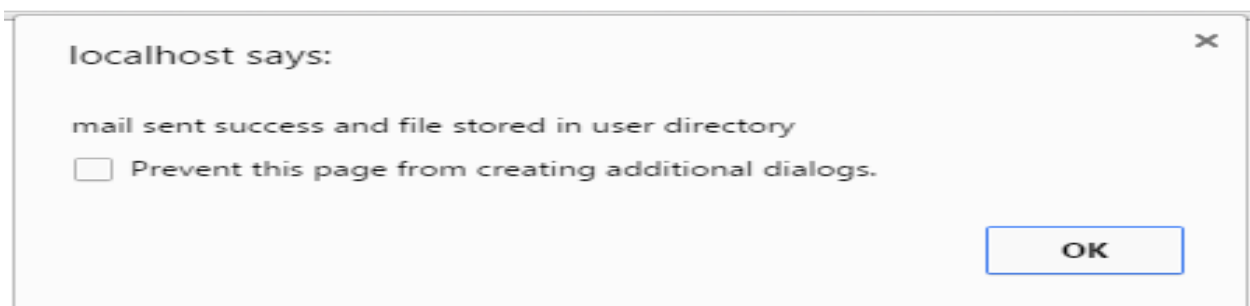


Fig.6: COMPOSE MAIL: In this mail user can upload files.



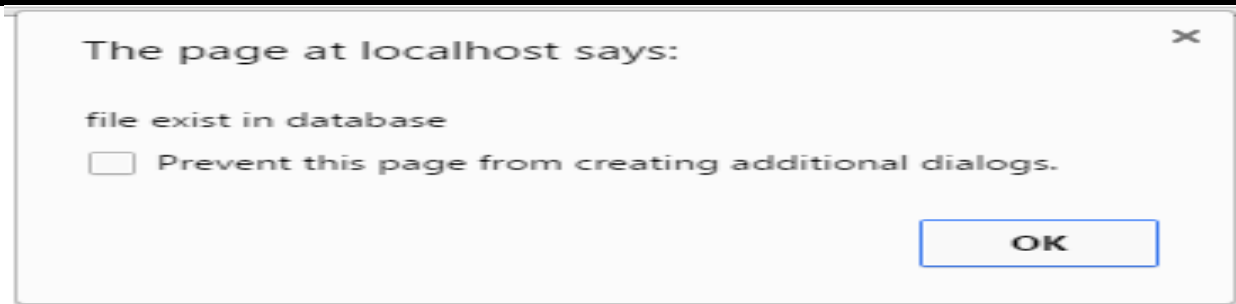


Fig.7: PATTERN MATCHING



Fig. 8: FORWARD FILES

VI. CONCLUSION & FUTURE WORK

By using Boyer-Moore algorithm, we can maintain the single file in the server. By maintaining single file we can save the memory space of the server. There may be replicated files be uploaded again and again so replication

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